

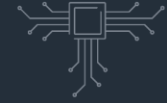
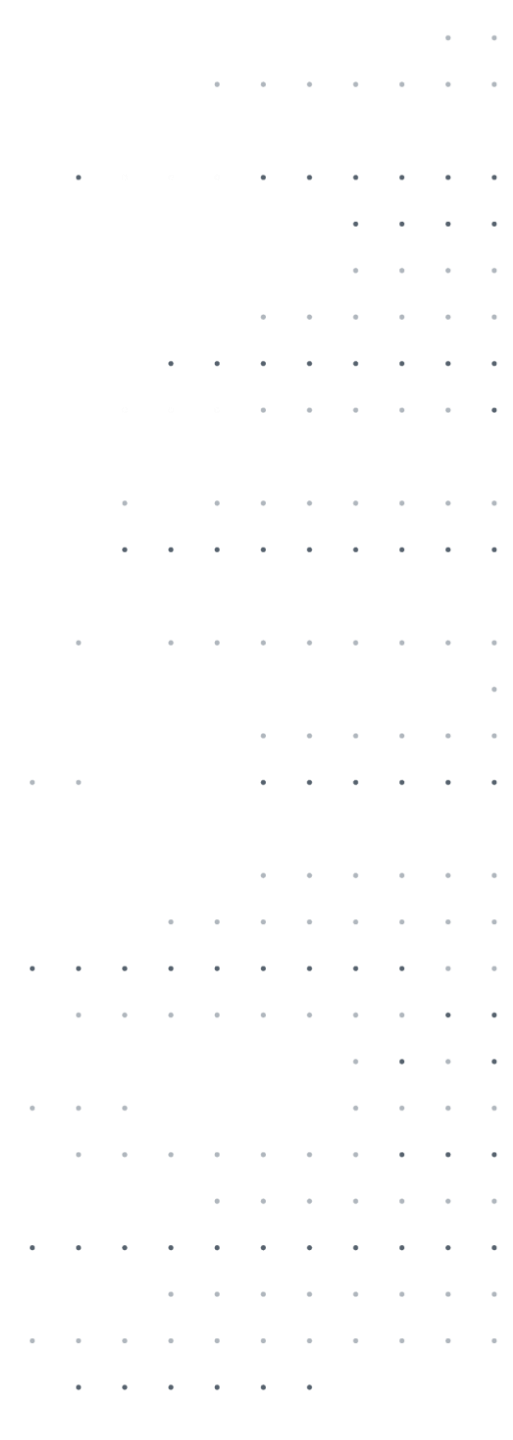
Challenges to gate-based quantum optimization algorithms for industrial use-cases



Quantum Computing
Research Leonardo

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Electronics



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Aerostructures

Combinatorial optimization: the antenna placement problem

Ising problems:

$$H_S = \sum_{i \neq j}^N W_{ij} s_i s_j + \sum_{i=1}^N A_i s_i \quad \text{with} \quad s_i \in \{-1, 1\}$$

Solution: combination of $S = \{s_i\}$ that minimizes H_S ?

Industrial problems:

1. Unstructured problems with few or no symmetries
2. Weighted graphs, Constraints (graphs become fully connected)
3. Focus is on the solution string

Countless applications: logistics, task scheduling, antenna placement, satellite mission planning, resource management, components on an aircraft...



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Use case in our work

M. Vandelli, A. Lignarolo, C. Cavazzoni, D. Dragoni, arXiv:2311.11621

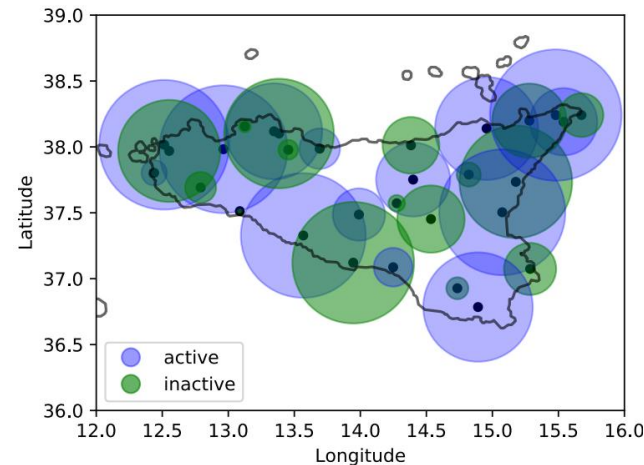
Optimization of a network of antennas operating in an emergency scenario:

- with scarce resources
- only a specified number of antennas can be simultaneously operated.

- Maximize the coverage of the signal generated by the antennas
- Minimize the overlaps between the signals
- Limited number of antennas

Antenna problem

Ising model



$$H(z; \xi, \lambda, N_t) = H_O(z) + H_C(z) + H_P(z)$$

with

$$H_O(z) = \sum_{i \neq j} z_i J_{ij} z_j$$

$$H_C(z) = -\xi \sum_i A_i z_i$$

$$H_P(z) = \lambda \left(\sum_{i \neq j} z_i z_j + \delta N_t \sum_i z_i \right)$$

$\lambda = 0 \rightarrow$ unconstrained
 $\lambda > 0 \rightarrow$ constrained



Emulating quantum algorithms on the *davinci-1* cluster

Two algorithms described by the same circuit:

$$|\beta, \gamma\rangle = U(\beta, \gamma) |+\rangle^{\otimes n} \quad \text{with} \quad U(\beta, \gamma) = \prod_{n=1}^p e^{-i\beta_n \sum_i X_i} e^{-i\gamma_n H_S}$$

QAA: linear Trotterization $\rightarrow \beta = \Delta \left(1 - \frac{k}{p}\right), \gamma = \Delta \frac{k}{p}$

QAOA: optimization in $2p$ -dimensions, $(\beta, \gamma) = \text{argmin} f(\beta, \gamma)$

[QAA: Fahri *et al.*, *arXiv:0001106*; QAOA: Fahri *et al.*, *arXiv:1411.4028*]



All the calculations performed on the proprietary *davinci-1* cluster equipped with AMD EPYC 7402 24-Core CPUs and NVIDIA A100 GPUs.

M. Vandelli, A. Lignarolo, C. Cavazzoni, D. Dragoni, *arXiv:2311.11621*

Total number of layers


$$p_{tot} = p \cdot N_{iter}$$

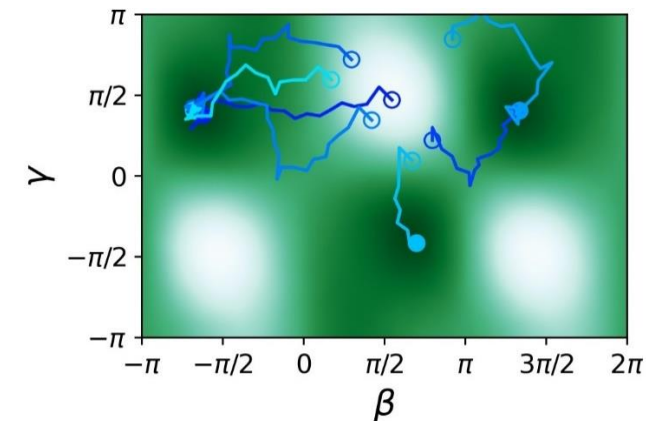
$$N_{iter}(QAA) = 1$$

$$N_{iter}(QAOA) = 50$$

$p \rightarrow$ Single-circuit depth

$p_{tot} \rightarrow$ Time complexity:
Total number of sequentially applied layers

- exact emulation based on  Qiskit
- MPI framework for global optimization using a **multi-walker optimization** initialized around an **INTERP** strategy
- COBYLA minimizer for local optimization



QAOA: 32 walkers x (8 CPU cores + 1 GPU)

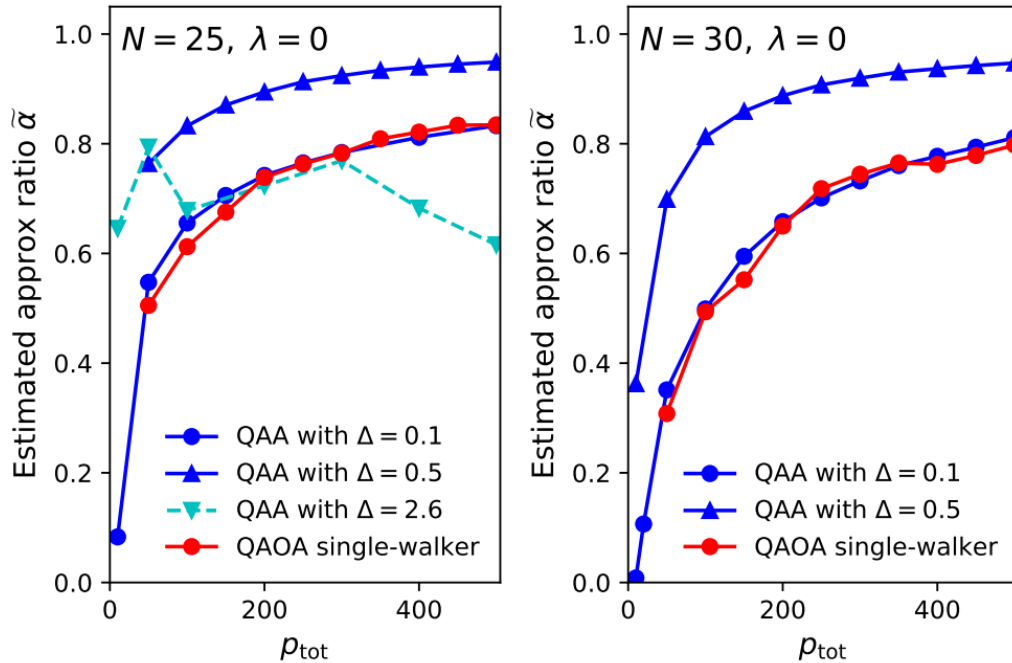
QAA & single-walker QAOA: 48 CPU cores + 1 GPU



Approximation ratio at large depths

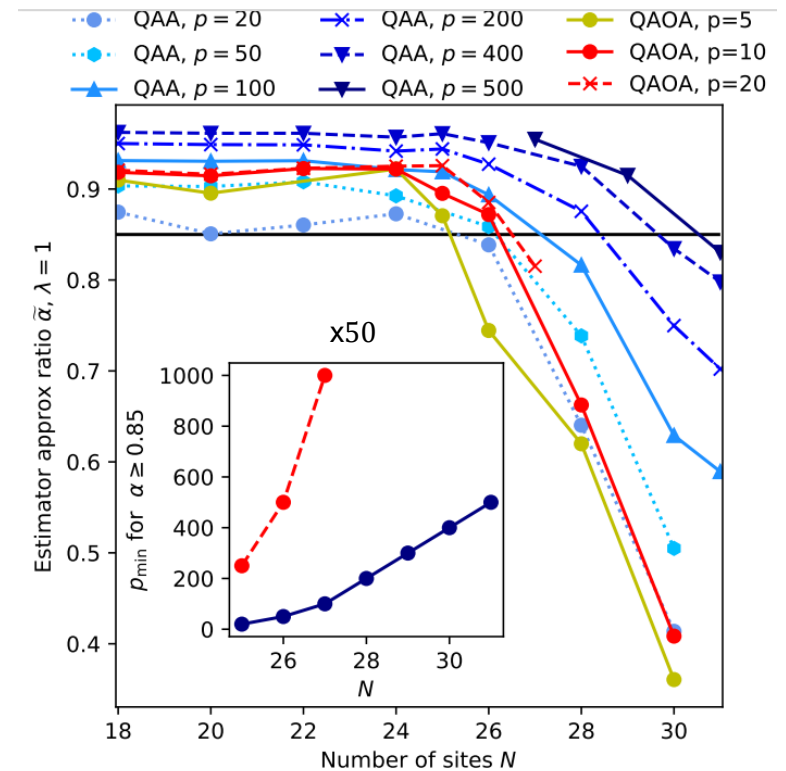
Approximation ratio: $\alpha = \langle H_\lambda \rangle / H_{min}$

Fixed shots estimator: $\tilde{\alpha} = \frac{\langle \sum_k H(z_k) \rangle / N_{meas}}{H_{min}}$



Approximation ratio as a function of the number of qubits:
Constrained case with $\lambda = 1$

$p_{min}(QAA) \sim 100 (N - 26)$
 $p_{min}(QAOA) \sim \exp(N)$ or $poly(N)$ (not linear)

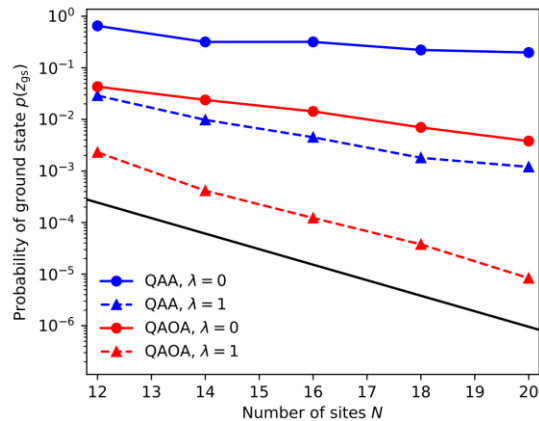


Generation of high-quality solutions

Probability of measuring the ground state

Emulation using the full statevector:

- QAOA has depth $p = 10$, while QAA has depth $p = 500$ (fixed computational time $p_{tot} = 500$ in both cases).
- In the constrained case ($\lambda = 1$), $p(z_{gs})$ is much lower.



Method	Measurements
QAA, $\lambda = 0$	1.07^N
QAOA, $\lambda = 0$	1.36^N
QAA, $\lambda = 1$	1.30^N
QAOA, $\lambda = 1$	1.67^N
Flat state	2.00^N

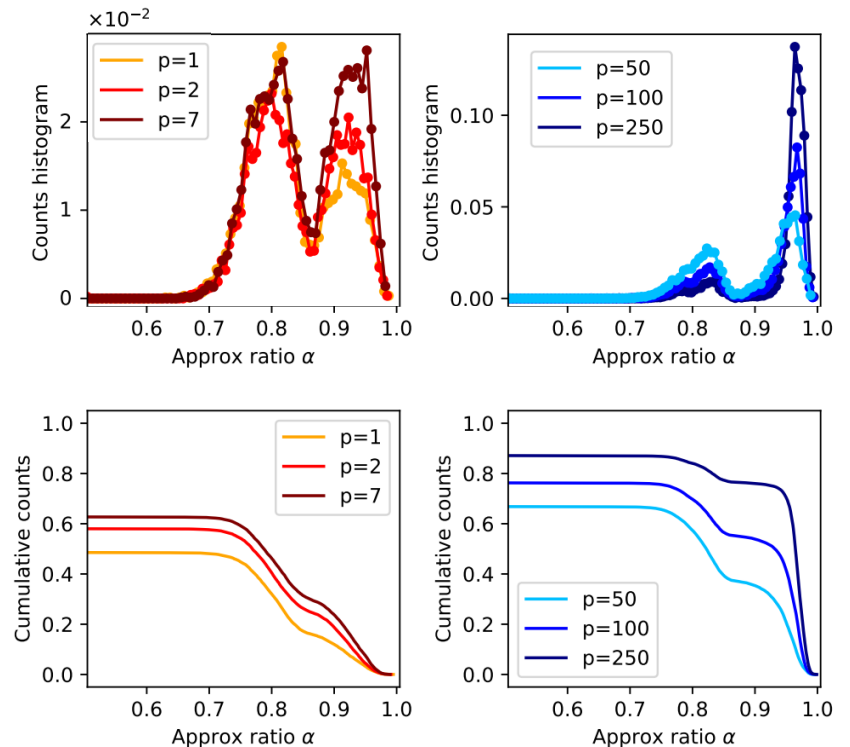
For this use case, the probability $p(z_{gs})$ decreases exponentially with the system size in all cases.

M. Vandelli, A. Lignarolo, C. Cavazzoni, D. Dragoni, arXiv:2311.11621

Generation of high-quality strings

$N = 30, N_{shots} = 10^4$, different values of p

- The cumulative counts estimate the probability of measuring a string with higher approx ratio.
- The probability is already large enough to guarantee a measurement of an acceptable string if we fix $\alpha \leq 0.95$

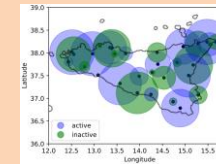


Summary & conclusions

Antenna placement problem: Unstructured problem with a weighted graph

$\lambda = 0$: sparse connectivity

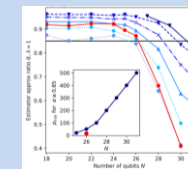
$\lambda > 0$: full connectivity



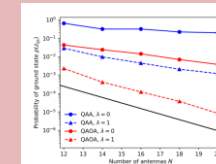
Approximation ratio with QAA and QAOA: rather robust against shot noise

Fixed α requires $p \sim N$ for QAA (and increases faster for QAOA)

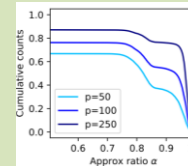
QAA: very deep circuits; **QAOA:** NP-hard optimization in p



Probability: for this prototypical unstructured problem, $p(z_{gs})$ decreases exponentially vs N for both QAA and QAOA



Cumulative distribution: These algorithms allow to measure good strings with high probability, if some tolerance is allowed



M. Vandelli, A. Lignarolo, C. Cavazzoni, D. Dragoni, arXiv:2311.11621





THANK YOU
FOR YOUR ATTENTION

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Academic vs practical optimization: our contribution

Quantum algorithms from an «academic» point of view:

1. Regular problems with symmetries
2. Unweighted graphs/integer parameters, sparse connectivity
3. Focus often on approximation ratio (cost, energy,...)

Practical point of view:

1. Unstructured problems with few or no symmetries
2. Weighted graphs, Constraints (graphs become fully connected)
3. Focus is on the solution string

Additionally, industrial problems are large → number of variables > 10^3 - 10^4

Comparison between 2 algorithms:

- 1) QAA
- 2) QAOA

Comparison between 2 different problems:

- 1) sparsely connected antennas ($\lambda = 0$, unconstrained)
- 2) Fully connected antennas ($\lambda = 1, N_t = \lfloor N/2 \rfloor$)

Metrics:

- 1) Approximation ratio α (single instance)
- 2) Probability of exact state $p(z_{gs})$
- 3) Cumulative probability $p(z | \alpha(z) \geq \alpha)$

Approximation ratio: $\alpha = \langle H_\lambda \rangle / H_{min}$

← Emulator of the full state-vector

Its estimator is $\tilde{\alpha} = \frac{(\sum_k H(z_k) / N_{meas})}{H_{min}}$

← Emulation with fixed number of shots

